

# On Dots in Boxes or Permutation Pattern Classes and Regular Languages

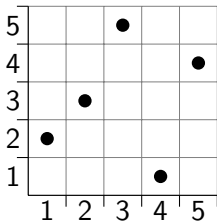
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- $a \in \Sigma$  is a regular language
- $\varepsilon$  is a regular language
- $\emptyset$  is a regular language
- if  $v$  and  $w$  are regular languages then so is  $v \cup w$
- if  $v$  and  $w$  are regular languages then so is  $vw$
- if  $v$  is a regular language then so is  $v^*$

# Basics – Rank Encoding



$$E(23514) = 22311$$

# Summary of Contributions

Proved that the following are regular languages under the rank encoding

- **Plus-(In)Decomposable Permutations**
- Minus-(In)Decomposable Permutations
- Direct Sum of Regular Classes
- $\sigma$ -Decomposable Permutations
- Simple Permutations

Proved that the following are not regular languages under the rank encoding

- Skew Sum of Regular Classes (in some cases)
- Inflation of Regular Classes (in most cases)
- Wreath Product of Regular Classes (in most cases)
- Wreath Closure of Regular Classes (in nearly all cases)

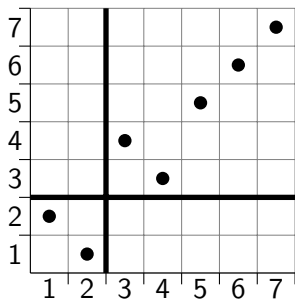
## Theorem (H. and Linton (2013))

*Let  $\mathcal{C}$  be a regular class. Then  $\mathcal{I}_P(\mathcal{C})$ , the set of all plus-indecomposable permutations of  $\mathcal{C}$ , is also regular under the rank encoding.*

## Corollary (H. and Linton (2013))

*Let  $\mathcal{C}$  be a regular class. Then  $\mathcal{D}_P(\mathcal{C})$ , the set of all plus-decomposable permutations of  $\mathcal{C}$ , is also regular under the rank encoding.*

# Regularity of Plus-(In)Decomposable Permutations



$$E(2143567) = 2121111$$

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# Regularity of Minus-(In)Decomposable Permutations

## Theorem (H. and Linton (2013))

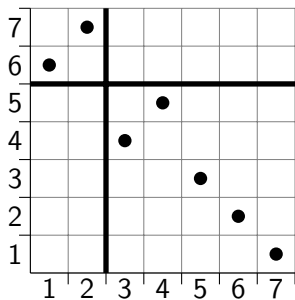
*Let  $\mathcal{C}$  be a regular class. Then  $\mathcal{D}_M(\mathcal{C})$ , the set of all minus-decomposable permutations of  $\mathcal{C}$ , is also regular under the rank encoding.*

## Corollary (H. and Linton (2013))

*Let  $\mathcal{C}$  be a regular class. Then  $\mathcal{I}_M(\mathcal{C})$  the set of all minus-indecomposable permutations of  $\mathcal{C}$ , is also regular under the rank encoding.*



# Regularity of Minus-(In)Decomposable Permutations



$$E(6745321) = 6644321$$

# Summary of Contributions

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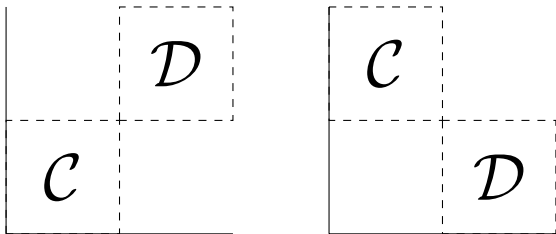
## Theorem

*Let  $\mathcal{C}$  and  $\mathcal{D}$  be two regular classes under the rank encoding. Then the direct sum  $\mathcal{E} = \mathcal{C} \oplus \mathcal{D}$  is a regular class under the rank encoding.*

## Theorem

*Let  $\mathcal{C}$  and  $\mathcal{D}$  be two regular classes under the rank encoding. Then the skew sum  $\mathcal{E} = \mathcal{C} \ominus \mathcal{D}$  is a regular class under the rank encoding if and only if  $\mathcal{D}$  is finite.*

# Regularity of the Direct and Skew Sum of Classes



# Summary of Contributions

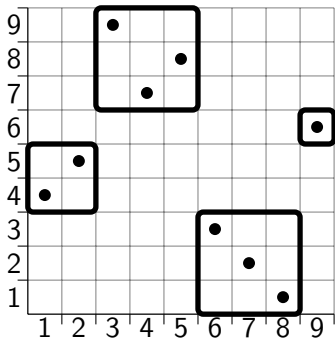
Proved that the following are regular languages under the rank encoding

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# Regularity of $\sigma$ -decomposable Permutations



## Theorem

Let  $E(\Omega_k)$  be the regular language of the rank encoded permutations with rank at most  $k$ . Let  $|\sigma| > 2$  be a simple permutation, then set  $\mathcal{D}_\sigma \subseteq \Omega_k$  of  $\sigma$ -decomposable permutations of  $\Omega_k$  is also regular under the rank encoding.

# Summary of Contributions

Proved that the following are regular languages under the rank encoding

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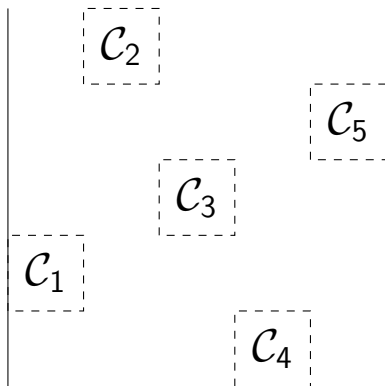
- Skew Sum of Regular Classes (in some cases)
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## Lemma (34)

*Let  $\sigma$  be any simple permutation of length  $n$ ,  $\sigma[\mathcal{C}_1, \dots, \mathcal{C}_n] = \mathcal{D}$  be an inflation of  $\sigma$  by the regular classes  $\mathcal{C}_i \subseteq \Omega_k$ . Then  $E(\mathcal{D})$  is a regular language under the rank encoding if and only if  $\mathcal{C}_j$  is finite when  $\sigma(j)$  is not a left to right maximum.*



# Regularity of the Inflation of Classes



# Summary of Contributions

Proved that the following are regular languages under the rank encoding

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- Minus-(In)Decomposable Permutations
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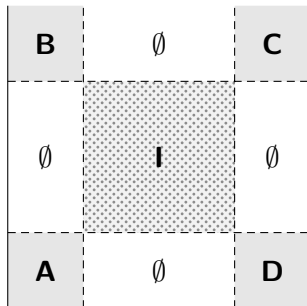
## Theorem

*The set of all non-simple permutations  $\mathcal{NS}_k$  of  $\Omega_k$  is regular under the rank encoding.*

## Corollary

*The set of simple permutations of  $\Omega_k$  is regular under the rank encoding.*

# Regularity of Non-Simple Permutations



# Regularity of Non-Simple Permutations

$$\begin{aligned} E_R(\mathcal{NS}_k) = & E_R(\Omega_k) \cap \\ & \left( \bigcup_{l=1}^{k-1} \mathcal{P}_l \bigcup_{m=l}^{k-1} m + E_R(\hat{\Omega}_{k-m}) \cup \bigcup_{j=1}^{k-1} j + E_R(\hat{\Omega}_{k-j}) \cup \right. \\ & \left. \bigcup_{a=2}^{k-1} \bigcup_{b=0}^{k-1-a} \mathcal{Q}_{a,b} \bigcup_{i=0}^{a-2} ((b+i) + E_R(\hat{\Omega}_{k-(b+i)}))^{(a-i)} \cup \right. \\ & \left. E_R(\Omega_k \setminus \{\varepsilon\}) E_R(\Omega_k \setminus \{\varepsilon\}) \right) \Sigma^* \end{aligned}$$

## Theorem

*Let  $\mathcal{C}$  be a permutation class. Let  $Si(\mathcal{C})$  be the set of all simple permutations of  $\mathcal{C}$ . If  $Si(E_k(\mathcal{C})) = Si(E_{k+1}(\mathcal{C})) = Si(E_{k+2}(\mathcal{C}))$  then  $Si(E_k(\mathcal{C}))$  is the set of words corresponding to all simple permutations of  $\mathcal{C}$ .*

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## Lemma (42)

*Let  $\mathcal{A}$  and  $\mathcal{B}$  be regular languages under the rank encoding then  $\mathcal{A} \wr \mathcal{B}$  is regular under the rank encoding if and only if  $\mathcal{B}$  is finite or  $\mathcal{A} \subseteq Av(21)$ .*



## Lemma (43)

*The wreath closure of a regular class  $\mathcal{A}$  containing the permutation 21 is not regular.*

## Corollary

*A wreath closed class  $\mathcal{A}$  is regular if and only if  $\mathcal{A}$  is finite or consists of ascending permutations.*

Is there a constructive way of finding the language of the basis of a geometrically griddable class?

# Thank you

Thanks!

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